## Introduction to Computer Security

## Discussion 10

Q1 Attack on TCP

(4 points)

Client (p	port $P_C$ )	Server (p
	SYN	
		,
	SYN-ACK	
	ACK	
	DATA A	·
		*
	DATA B	
	DATA C	
		,

- Server (port  $P_S$ )
  - 1. Client sends initial SYN with sequence number X (usually random).
  - Server sends SYN-ACK with sequence number Y (also usually random) and ACK X + 1.
  - 3. Client sends ACK with sequence number X+1 and ACK Y+1.
  - 4. Client sends DATA A of length  $L_A$  with sequence number X+1 and ACK Y+1.
  - 5. Server sends DATA B of length  $L_B$  with sequence number Y+1 and ACK  $X+1+L_A$ .
  - 6. Client sends DATA C of length  $L_C$  with sequence number  $X+1+L_A$  and ACK  $Y+1+L_B$ .
  - Data exchange continues until both sides are done sending data.
- Q1.1 (1 point) Assume that the next transmission in this connection will be DATA D from the server to the client. What will this packet look like?

Sequence Number:		ACK:	
Source Port:	$P_S$	Destination Port:	$P_C$
Length:	$L_D$	Flags:	ACK

Q1.2	(1 point) You should be familiar with the concept and capabilities of a <i>man-in-the-middle</i> as an attacker who <b>can observe</b> and <b>can modify</b> traffic. There are two other types of relevant attackers in this scenario:
	<ol> <li>On-path attacker: can observe traffic but cannot modify it.</li> <li>Off-path attacker: cannot observe traffic and cannot modify it.</li> </ol>
	Carol is an <i>on-path</i> attacker. Can Carol do anything malicious to the connection? If so, what can she do?
Q1.3	(1 point) David is an <i>off-path</i> attacker. Can David do anything malicious to the connection? If so, what can he do?
Q1.4	(1 point) The client starts getting responses from the server that don't make any sense. Inferring that David is attempting to hijack the connection, the client then immediately sends the server a <b>RST</b> packet, which terminates the ongoing connection. David wants to impersonate the client by establishing a new connection. How would he go about doing this?

(Question 1 continued...)

Q2 TLS Protocol Details (4 points)

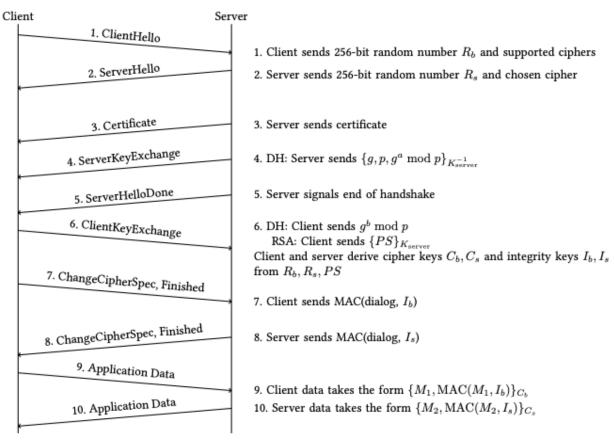


Figure 1: TLS 1.2 Key Exchange

Q2.1	(1 point) What is the purpose of the <i>client random</i> and <i>server random</i> fields?
Q2.2	(1 point)
	ClientHello and ServerHello are not encrypted or authenticated. Explain why a man-in-the-middle cannot exploit this. (Consider both the Diffie-Hellman and RSA case.)

	cipher key $C$ and one integrity key $I$ . Why might this be a bad idea?
Q2.4	4 (1 point)
	The protocol given above is a simplified form of what actually happens. After step 8 (ChangeCipherSpec), the protocol as described above is still vulnerable. What is the vulnerability and how could you fix this?

(Question 2 continued...)

Q3 TLS Threats (3 points)

An attacker is trying to attack the company Boogle and its users. Assume that users always visit Boogle's website with an HTTPS connection, using ephemeral Diffie-Hellman. You should also assume that Boogle does not use certificate pinning. The attacker may have one of three possible goals:

- 1. Impersonate the Boogle web server to a user
- 2. Discover some of the plaintext of data sent during a past connection between a user and Boogle's website
- 3. Replay data that a user previously sent to the Boogle server over a prior HTTPS connection

For each of the following scenarios, describe if and how the attacker can achieve each goal. Q3.1 (1 point) The attacker obtains a copy of Boogle's certificate.

Q3.2	(1 point) The attacker obtains the private key of a certificate authority trusted by users of Boogle.

Q3.3 (1 point) The attacker obtains the private key corresponding to an old certificate used by Boogle's server during a past connection between a victim and Boogle's server. Assume that this old certificate has been revoked and is no longer valid. Note that the attacker does not have the private key corresponding to current certificate.