CS161 Fall 2025

Introduction to Computer Security

Discussion 7

Q1 Boogle (5 points)

Boogle is a social networking website that's looking into expanding into other domains. Namely, they recently started a map service to try their hand at fusing that with social media. The URL for the main website is https://www.boogle.com, and they want to host the map service at https://maps.boogle.com.

- Q1.1 (1 point) For each of the following webpages, determine whether the webpage has the same origin as http://boogle.com/index.html, and provide a brief justification.
 - i. https://boogle.com/index.html
 - ii. http://maps.boogle.com
 - iii. http://boogle.com/home.html
 - iv. http://maps.boogle.com:8080

Solution:

- i. False. https://boogle.com/index.html do not have the same origin, since their protocols (https) and (http) are different.
- ii. False. http://boogle.com/index.html do not have the same origin, since their domains (maps.boogle.com) and (boogle.com) are different. The same-origin policy performs string matching on the protocol, domain, and port.
- iii. True. The paths are not checked in the same-origin policy.
- iv. False. http://maps.boogle.com:8080 and http://boogle.com/index.html do not have the same origin, because their ports (8080) and (80) are different. Note that if the port is not specified, the port defaults to 80 for http and 443 for https.
- Q1.2 (1 point) Describe how to make a cookie that will be sent to only Boogle's map website and its subdomains.

Solution: Set the domain parameter of the cookie to .maps.boogle.com

Q1.3 (1 point) How can Boogle ensure that cookies are only transmitted encrypted so eavesdroppers on the network can't trivially learn the contents of the cookies?

Solution: Set the secure flag on each cookie.

Q1.4 (1 point) Boogle wants to be able to host websites for users on their servers. They decide to host each user's website at https://[username].boogle.com. Why might this not be a good idea?

Solution: A malicious user could set cookies that would be sent to other users' sites as well as the entire .boogle.com domain. Also, any cookies meant for boogle.com will go to the malicious user.



(Question 1 continued...)

Q1.5 (1 point) Propose an alternate scheme so that Boogle can still host other users websites with less risk, and explain why this scheme is better.

Note: It is okay if the user sites interfere with each other, as long as they cannot affect official Boogle websites.

Solution: Boogle should create a new domain exclusively for user hosted content, like https://lisername].boogleusercontent.com. This way, user sites cannot set cookies that will affect all boogle domains due to the cookie setting policy. This is known as a cookie tossing attack, and is one of the reasons why github hosts user sites on github.io instead of github.com (see https://blog.github.com/2013-04-09-yummy-cookies-across-domains/).

Q2 Session Fixation (2 points)

A *session cookie* is used by most websites in order to manage user logins. When the user logs in, the server sends a randomly-generated session cookie to the user's browser. The server also stores the cookie value in a database along with the corresponding username.

The user's browser sends the session cookie to the server whenever the user loads any page on the site. The server then looks the session cookie up in the database and retrieves the corresponding username. Using this, the server can know which user is logged in.

Some web application frameworks allow cookies to be set by the URL. For example, visiting the URL

http://foobar.edu/page.html?sessionid=42

will result in the server setting the sessionid cookie to the value "42".

Q2.1 (1 point) Can you spot an attack on this scheme?

Solution: The main attack is known as session fixation. Say the attacker establishes a session with foobar.edu, receives a session ID of 42, and then tricks the victim into visiting http://foobar.edu/browse.html?sessionid=42 (maybe through an img tag). The victim is now browsing foobar.edu with the attacker's account. Depending on the application, this could have serious implications. For example, the attacker could trick the victim to pay his bills instead of the victim's (as intended).

Another possibility is for the attacker to fix the session ID and then send the user a link to the log-in page. Depending on how the application is coded, it might so happen that the application allows the user to log-in but reuses the previous (attacker-set) session ID. For example, if the victim types in his username and password at http://foobar.edu/login.html?sessionid=42, then the session ID 42 would be bound to his identity. In such a scenario, the attacker could impersonate the victim on the site. This is uncommon nowadays, as most login pages reset the session ID to a new random value instead of reusing an old one.

Q2.2 (1 point) Suppose the problem you spotted has been fixed as follows: foobar.edu now establishes new sessions with session IDs based on a hash of the tuple (username, time of connection). Is this secure? If not, what would be a better approach?

Solution: The proposed fix is not secure since it solves the wrong problem - it doesn't fix either issue. In fact, it makes things weaker by significantly reducing the *entropy* of the session cookie.

The correct fix is for the server to generate cookie values afresh, rather than setting them based on the session ID provided via URL parameters. Also, the server shouldn't allow cookies to be set by the URL. This makes the attackers job more difficult as they have to do some form of XSS in order to manipulate the client's cookie vs. just clicking on a link.

In a CSRF attack, a malicious user is able to take action on behalf of the victim. Consider the following example. Mallory posts the following in a comment on a chat forum:

```
<img src="http://patsy-bank.com/transfer?amt=1000\&to=mallory"/>
```

Of course, Patsy-Bank won't let just anyone request a transaction on behalf of any given account name. Users first need to authenticate with a password. However, once a user has authenticated, Patsy-Bank associates their session ID with an authenticated session state.

Q3.1 (1 point) Explain what could happen when Alice visits the chat forum and views Mallory's comment.

Solution: The img tag embedded in the form causes the browser to make a request to http://patsy-bank.com/transfer?amt=1000\&to=mallory

with Patsy-Bank's cookie. If Alice was previously logged in (and didn't log out), Patsy-Bank might assume Alice is authorizing a transfer of 1000 USD to Mallory.

Q3.2 (1 point) Patsy-Bank decides to check that the **Referer** header contains patsy-bank.com. Will the attack above work? Why or why not?

Solution: In most cases, it will solve the problem since the **Referer** header will contain the blog's URL instead of patsy-bank.com.

However, not all browsers send the **Referer** header, and even when they do, not all requests include it.

Q3.3 (1 point) Describe one way Mallory can modify her attack to always get around this check.

Solution: She can have the link go to a URL under Mallory's control which contains patsybank.com such as

(patsy-bank.com).attacker.com

or

attacker.com/attack?dummy=patsy-bank.com.

Then this page can redirect to the original malicious link. Now the **Referer** header will pass the check.

Another solution, is if the Patsy-Bank has a so-called "open redirect"

http://patsy-bank.com/redirect?to=url,

the referrer for the redirected request will be

http://patsy-bank.com/redirect?to=....

An attacker can abuse this functionality by causing a victim's browser to fetch a URL like http://patsy-bank.com/redirect?to=http://patsy-bank.com/transfer..., and from patsy-bank.com's perspective, it will see a subsequent request for http://patsy-bank.com/transfer... that indeed has a Referer from patsy-bank.com.

Q3.4 (1 point) Recall that the **Referer** header provides the full URL. HTTP additionally offers an **Origin** header which acts the same as the **Referer** but only includes the website domain, not the entire URL. Why might the **Origin** header be preferred?

Solution: Leaking the entire URL can be a violation of privacy against users. As an example, consider Alice transferred money by visiting

http://patsy-bank.com/transfer?amt=1000\&to=bob

and subsequently went to a website under an attacker's control - now the attacker has learned the exact amount of money Alice sent and to who. The Origin header would only leak that Alice was at the patsy-bank.com.

As a sidenote not directly related to the question, the Origin is a very useful way to solve the CSRF problem since it makes it much easier for multiple, trusted sites to make some action. For example, Patsy-Bank might trust

http://www.trustedcreditcardcompany.com

to directly transfer money from a user's account. This is a use-case that the CSRF token-based solution doesn't support cleanly.

Q3.5 (1 point) Almost all browsers support an additional cookie field SameSite. When SameSite=strict, the browser will only send the cookie if the requested domain and origin domain correspond to the cookie's domain. Which CSRF attacks will this stop? Which ones won't it stop? Give one big drawback of setting SameSite=strict.

Solution: It stops almost all CSRF attacks, except those involving open redirects from the website in question or if the website itself has an XSS vulnerability.

However, setting SameSite=strict can greatly limit functionality since any external links that require a user to be logged in won't work. For instance, consider a friend sends you a Facebook link via email, clicking on that link will require you to sign in again since your session cookie wasn't sent with the request.